# On the evolution of the expected gain of a greedy action in the bandit problem 

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#### Abstract

The K-armed bandit problem is a well-known formalization of the exploration versus exploitation dilemma. In a K-armed bandit problem, a player is confronted with a gambling machine with $K$ arms where each arm is associated to an unknown gain distribution and the goal is to maximize the sum of the rewards. One of the simplest policies for this bandit problem is the greedy policy which keeps a gain estimation of the arms and at each round greedily chooses the arm which, on average, performed the best so far. This paper defines and gives an analytical definition of the expected gain of a greedy action $\mu_{g}$ and studies its evolution over the time. If the gambling machine has two arms, we derive an optimal exploration algorithm for the two arms case and we show that the evolution of $\mu_{g}$ under an exploitation greedy policy is upper-bounded. Otherwise if the player is confronted with a gambling machine with more than two arms we show experimentally that the $\mu_{g}$ evolution is much more complex and that the value of $\mu_{g}$ can decrease. Our results confirm analytically and experimentally that exploitation before acquiring enough knowledge on the arms is a bad practice.


In this paper, we study the well-known K-armed bandit problem, first introduced by Robbins [6]. In this learning problem, a casino player has to decide which arm of a K-slot machine to pull to maximize the total gain in a series of rounds. Each of the K arms of the slot machine returns a reward which is randomly distributed and unknown to the player. The player has to define a sequential selection policy on the basis of limited knowledge about the reward distributions which derive exclusively from previous results.

The K-armed bandit problem is a classical instance of situations involving a exploration/exploitation trade-off. An example of such a situation is the design of clinical trials to assess and compare a set of new medical treatments [4]. Here, the goal is to determine the best treatment while minimizing the inconveniences for the patients. Another example is the business problem of selecting the best supplier on the basis of incomplete information [2].

Several approaches have been proposed in literature to deal with the K-armed bandit problem. In this paper, we will focus on semi uniform algorithms [7]. A semi uniform algorithm is characterized by the alternation of two working modes, namely the exploration mode and the exploitation mode. The simplest semi uniform algorithm is the greedy policy in which the player keeps an updated
estimation of the gains of the arms and at each round greedily chooses the arm which, on average, performed the best so far. If the player adopts the greedy policy, we say that he adopts a pure exploitation policy. This means that he uses current knowledge to select the seemingly best arm without reserving any time to explore what seems to be inferior arms. Exploitation is the right thing to do to maximize the short term gain of an action and exploration is the right thing to do to maximize gain in long term.

A possible alternative to pure exploitation is the $\epsilon$-greedy policy [8] which preserves a fixed fraction of the time, e.g., quantified by a parameter $\epsilon$, to perform random uniform selection. A variant of the $\epsilon$-greedy is the $\epsilon$-decreasing-greedy policy where the exploration rate is set initially to a high value and then gradually decreases. For specific conditions on the initial parameters of $\epsilon$-decreasinggreedy, Auer et al. in [1] found an upper decreasing bound on the probability of selecting a sub-optimal arm. This means that after enough rounds, a player adopting the $\epsilon$-decreasing-greedy policy has a high probability of always playing the best arm.

The probability that a greedy action selects the true best arm is well-known in simulation literature as the probability of correct selection (PCS) [5]. In that case the issue is to decide how many simulation trials should be conducted if we want to have a certain guarantee that the correct selection will be accomplished. In [3], we propose to use the PCS notion as a founding principle of a sequential strategy and as a measure of the effectiveness of an exploration step. In this paper, we extend existing work by (i) extending the notion of PCS to the notion of expected greedy reward and (ii) studying the behavior of the evolution in the time of this expected greedy reward.

This paper is structured as follows : Section 1 defines formally the bandit problem and the greedy action. An analytical definition of the expected greedy reward is given in Section 2 and the behavior of its evolution is studied in section 3.

## 1 The bandit problem

This section formally defines the K-armed bandit problem and introduces the notations used in this paper.

A K-armed bandit problem can be modeled by a set $\mathbf{z}_{\#}=\left\{\mathbf{z}_{k}\right\}, k=1, \ldots, K$ of $K$ random rewards ${ }^{1} \mathbf{z}_{k}$ with mean $\mu_{k}$ and standard deviation $\sigma_{k}$. Suppose that the goal of the player is to maximize the collected rewards. Once fixed the duration of the game to $H$ rounds, at each round $l, l=1, \ldots, H$ the player is expected to select an arm out of the $K$ alternatives. Let $\mathbf{N}(l)=\left[\mathbf{n}_{1}(l), \ldots, \mathbf{n}_{K}(l)\right]$ be a counting vector whose $k$ th term denotes the number of times that the $k$ th arm was selected during the $l$ first rounds and $\mathbf{Z}_{k}(l)=\left[\mathbf{z}_{k}^{1}, \mathbf{z}_{k}^{2}, \mathbf{z}_{k}^{3}, \ldots, \mathbf{z}_{k}^{\mathbf{n}_{k}(l)}\right]$ be the vector of identically and independently distributed observed rewards of the arm $k$ up to time $l$.

[^0]Based on the samples in the set $\left\{\mathbf{Z}_{k}(l)\right\}, k=1, \ldots, K$ and following a policy, the player selects iteratively one arm $\widehat{\mathbf{k}}$ at each round $l$. We define (i) the observed state $\widehat{\mathbf{s}}(l) \in \widehat{\mathcal{S}}$ of the game at the $l$ th round as the whole set of observations $\left\{\mathbf{Z}_{k}(l)\right\}, k=1, \ldots, K$ and (ii) the policy as a function $\widehat{\pi}: \widehat{\mathcal{S}} \rightarrow\{1, \ldots, K\}$ which returns for each state $\widehat{\mathbf{s}}$ the $\operatorname{arm} \widehat{\mathbf{k}}$.

Two common adopted strategies are the pure random policy and the greedy policy. The pure random policy neglects any partial information about the state at time $l$ and returns a selection

$$
\widehat{\mathbf{k}}_{r}=\widehat{\pi}(\widehat{\mathbf{s}}(l)) \sim \operatorname{Uni}(1 / K, \ldots, 1 / K)
$$

sampled according to an uniform distribution.
The greedy policy uses instead the information contained in $\widehat{\mathbf{s}}(l)$ and returns

$$
\widehat{\mathbf{k}}_{g}=\arg \max \widehat{\boldsymbol{\mu}}_{k}(l)
$$

where $\widehat{\boldsymbol{\mu}}_{k}(l)$ is the sampled average of the vector $\mathbf{Z}_{k}(l)$.

## 2 An analytical definition of the expected greedy reward

This section introduces and gives an analytical definition of the expected gain of a greedy exploitation action $\mu_{g}$. Since $\mathbf{Z}_{k}(l)$ is a realization of a random vector, the state $\widehat{\mathbf{s}}$ and the output $\widehat{\mathbf{k}}=\widehat{\pi}(\widehat{\mathbf{s}})$ of the policy are random realizations too. At round $l$ it is then possible to associate to a policy the expected gain $\mathbb{E}\left[\mu_{\hat{\mathbf{k}}}\right]$ which quantifies the gain caused by the policy adoption at round $l$. Note that $\mu_{\widehat{\mathbf{k}}}$ is a random variable which denotes the mean of the arm $\widehat{\mathbf{k}}$ selected at the $l \mathrm{th}$ step.

Let

$$
\begin{equation*}
P_{\bar{k}}(l)=\operatorname{Prob}\left\{\arg \max \widehat{\boldsymbol{\mu}}_{k}(l)=\bar{k}\right\} \tag{1}
\end{equation*}
$$

be the probability that a greedy algorithm selects arm $\bar{k}$ at the $l$ th step. The expected gain of a greedy exploitation action at round $l$ is then defined as follows

$$
\begin{equation*}
\mathbb{E}\left[\mu_{\widehat{\mathbf{k}}}\right]=\sum_{k=1}^{K} P_{k}(l) \cdot \mu_{k}=\mu_{g} . \tag{2}
\end{equation*}
$$

Note that $\mu_{g}<\mu_{k^{*}}$ where $k^{*}=\arg \max _{k} \mu_{k}$. Theorem_1 will show that an analytical expression of the probability $P_{\bar{k}}$ of selecting the $\bar{k}$ machine in a greedy strategy can be derived in case of a bandit problem with normally distributed arms.

## Theorem 1.

Let $\mathbf{z}_{\#}=\left\{\mathbf{z}_{1}, \ldots, \mathbf{z}_{K}\right\}$ be a set of $K>1$ normal reward distributions $\mathbf{z}_{k} \sim$ $\mathcal{N}\left[\mu_{k}, \sigma_{k}\right]$ with mean $\mu_{k}$ and standard deviation $\sigma_{k}$. If the selection policy is greedy then the probability of selecting $\mathbf{z}_{\bar{k}}$ is

$$
P_{\bar{k}}=\operatorname{Prob}\left\{\widehat{\mathbf{r}}_{1}>0, \ldots, \widehat{\mathbf{r}}_{\bar{k}-1}>0, \widehat{\mathbf{r}}_{\bar{k}+1}>0, \ldots, \widehat{\mathbf{r}}_{K}>0\right\}
$$

where $\left(\widehat{\mathbf{r}}_{1}, \ldots, \widehat{\mathbf{r}}_{\bar{k}-1}, \widehat{\mathbf{r}}_{\bar{k}+1}, \ldots, \widehat{\mathbf{r}}_{K}\right)^{T}$ follows a multivariate normal distribution

$$
\left(\widehat{\mathbf{r}}_{1}, \ldots, \widehat{\mathbf{r}}_{\bar{k}-1}, \widehat{\mathbf{r}}_{\bar{k}+1}, \ldots, \widehat{\mathbf{r}}_{K}\right)^{T} \sim \mathcal{N}[\Gamma, \Sigma]
$$

with mean

$$
\Gamma=\left(\begin{array}{c}
\mu_{\bar{k}}-\mu_{1} \\
\vdots \\
\mu_{\bar{k}}-\mu_{\bar{k}-1} \\
\mu_{\bar{k}}-\mu_{\bar{k}+1} \\
\vdots \\
\mu_{\bar{k}}-\mu_{K}
\end{array}\right)
$$

and covariance matrix $\Sigma$

$$
\Sigma=\left(\begin{array}{cccccc}
\frac{\sigma_{k}^{2}}{n_{\bar{k}}}+\frac{\sigma_{1}^{2}}{n_{1}} & \cdots & \frac{\sigma_{k}^{2}}{n_{\bar{k}}} & \frac{\sigma_{k}^{2}}{n_{\bar{k}}} & \cdots & \frac{\sigma_{k}^{2}}{n_{\bar{k}}} \\
\vdots & \ddots & \vdots & \vdots & & \vdots \\
\frac{\sigma_{\bar{k}}^{2}}{n_{\bar{k}}} & \cdots & \frac{\sigma_{\bar{k}}^{2}}{n_{\bar{k}}}+\frac{\sigma_{\bar{k}-1}^{2}}{n_{\bar{k}-1}} & \frac{\sigma_{\bar{k}}^{2}}{n_{\bar{k}}} & \cdots & \frac{\sigma_{\bar{k}}^{2}}{n_{\bar{k}}} \\
\frac{\sigma_{k}^{2}}{n_{\bar{k}}} & \cdots & \frac{\sigma_{k}^{2}}{n_{\bar{k}}} & \frac{\sigma_{k}^{2}}{n_{\bar{k}}}+\frac{\sigma_{k+1}^{2}}{n_{\bar{k}+1}} & \cdots & \frac{\sigma_{k}^{2}}{n_{\bar{k}}} \\
\vdots & & \vdots & \vdots & \ddots & \vdots \\
\frac{\sigma_{k}^{2}}{n_{\bar{k}}} & \cdots & \frac{\sigma_{k}^{2}}{n_{\bar{k}}} & \frac{\sigma_{k}^{2}}{n_{\bar{k}}} & \cdots & \frac{\sigma_{k}^{2}}{n_{\bar{k}}}+\frac{\sigma_{K}^{2}}{n_{K}}
\end{array}\right)
$$

where $n_{k}$ is the number of observations of $\mathbf{z}_{k}$.
Proof. According to (1), $P_{\bar{k}}$ is the probability that $\widehat{\boldsymbol{\mu}}_{\bar{k}}$ is the maximum in $\left\{\widehat{\boldsymbol{\mu}}_{1}, \ldots, \widehat{\boldsymbol{\mu}}_{K}\right\}:$

$$
\begin{aligned}
P_{\bar{k}} & =\operatorname{Prob}\left\{\widehat{\mathbf{k}}_{g}=\bar{k}\right\} \\
& =\operatorname{Prob}\left\{\bar{k}=\arg \max _{k \in[1 \ldots K]}\left\{\widehat{\boldsymbol{\mu}}_{k}\right\}\right\} \\
& =\operatorname{Prob}\left\{\widehat{\boldsymbol{\mu}}_{\bar{k}}>\widehat{\boldsymbol{\mu}}_{1}, \ldots, \widehat{\boldsymbol{\mu}}_{\bar{k}}>\widehat{\boldsymbol{\mu}}_{\bar{k}-1}, \widehat{\boldsymbol{\mu}}_{\bar{k}}>\widehat{\boldsymbol{\mu}}_{\bar{k}+1}, \ldots, \widehat{\boldsymbol{\mu}}_{\bar{k}}>\widehat{\boldsymbol{\mu}}_{K}\right\} \\
& =\operatorname{Prob}\left\{\widehat{\mathbf{r}}_{1}>0, \ldots, \widehat{\mathbf{r}}_{\bar{k}-1}>0, \widehat{\mathbf{r}}_{\bar{k}+1}>0, \ldots, \widehat{\mathbf{r}}_{K}>0\right\},
\end{aligned}
$$

where $\widehat{\mathbf{r}}_{k}=\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{k}$. It follows that $P_{\bar{k}}$ denotes also the probability that all the components of the vector $\left(\widehat{\mathbf{r}}_{1}, \ldots, \widehat{\mathbf{r}}_{\bar{k}-1}, \widehat{\mathbf{r}}_{\bar{k}+1}, \ldots, \widehat{\mathbf{r}}_{K}\right)^{T}$ are positive. Under the assumption of Gaussianity, this vector is a multivariate normal random variable

$$
\left(\begin{array}{c}
\widehat{\mathbf{r}}_{1} \\
\vdots \\
\widehat{\mathbf{r}}_{\bar{k}-1} \\
\widehat{\mathbf{r}}_{\bar{k}+1} \\
\vdots \\
\widehat{\mathbf{r}}_{K}
\end{array}\right)=\left(\begin{array}{c}
\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{1} \\
\vdots \\
\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{\bar{k}-1} \\
\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{\bar{k}+1} \\
\vdots \\
\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{K}
\end{array}\right) \sim \mathcal{N}[\Gamma, \Sigma]
$$

with mean vector

$$
\Gamma=\left(\begin{array}{c}
\mu_{\bar{k}}-\mu_{1} \\
\vdots \\
\mu_{\bar{k}}-\mu_{\bar{k}-1} \\
\mu_{\bar{k}}-\mu_{\bar{k}+1} \\
\vdots \\
\mu_{\bar{k}}-\mu_{K}
\end{array}\right)
$$

and covariance matrix

Now, since $\widehat{\boldsymbol{\mu}}_{i}$ and $\widehat{\boldsymbol{\mu}}_{j}$ are independent for $i \neq j$

$$
\sigma_{\widehat{\mathbf{r}}_{j}, \widehat{\mathbf{r}}_{j}}^{2}=\operatorname{Var}\left(\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{j}\right)=\frac{\sigma_{\bar{k}}^{2}}{n_{\bar{k}}}+\frac{\sigma_{j}^{2}}{n_{j}}
$$

it follows that

$$
\begin{aligned}
\sigma_{\widehat{\mathbf{r}}_{i}, \widehat{\mathbf{r}}_{j}}^{2}=\sigma_{\widehat{\mathbf{r}}_{j}, \widehat{\mathbf{r}}_{i}}^{2} & =\operatorname{cov}\left[\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{i}, \widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{j}\right] \\
& =\mathbb{E}\left[\left(\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{i}-\mathbb{E}\left[\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{i}\right]\right) \cdot\left(\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{j}-\mathbb{E}\left[\widehat{\boldsymbol{\mu}}_{\bar{k}}-\widehat{\boldsymbol{\mu}}_{j}\right]\right)\right] \\
& =\mathbb{E}\left[\left(\widehat{\boldsymbol{\mu}}_{\bar{k}}\right)^{2}\right]-\mu_{\bar{k}}^{2} \\
& =\frac{\sigma_{\bar{k}}^{2}}{n_{\bar{k}}} .
\end{aligned}
$$

## 3 The evolution of the expected greedy reward

A major difficulty of the bandit problem is the dynamic and multivariate nature of the terms involved in the definition of optimal policy. For this reason, it is important to study how the expected reward of a greedy action $\mu_{g}(2)$ changes with time.

Our analysis will be made in two parts. In the first one we assume that $\mathbf{z}_{\#}$ contains only two arms where $k^{*}$ is the index of the best arm and $\bar{k}^{*}$ is the other index. This is the simplest case and some analytical results can be derived about the evolution in time of $\mu_{g}$. The second part will analyse the most complex case, i.e. when $K>2$.

### 3.1 The evolution of $\mu_{g}$ in the case where $K=2$

First, we assume that the bandit machine has two arms. Theorem 2 shows that if $K=2$ then testing $\mathbf{z}_{1}$ or $\mathbf{z}_{2}$ at round $l$ will, in both cases, improve the value of $\mu_{g}$ at round $l+1$. Theorem 3 proposes an optimal exploration policy which maximizes $\mu_{g}$ at round $l+1$. Note that by definition, $\mu_{g}<\mu_{k^{*}}$. Theorem 4 shows that the evolution of $\mu_{g}$ of any strategy which always tests the same arm is upper-bounded by a value $u b^{\infty}$ which is smaller than $\mu_{k^{*}}$. This is an argument in favor of exploration.

Theorem 2. Let $K=2$ and $\mu_{g}^{l+1}(k)$ be the next expected reward of a greedy action when arm $k$ is tested at round $l$. Then

$$
\forall k \in\{1,2\}, \mu_{g}^{l+1}(k)>\mu_{g} .
$$

Proof. $\forall k \in\{1,2\}, \mu_{g}^{l+1}(k)>\mu_{g}$ if and only if $\forall k \in\{1,2\}, P_{k^{*}}^{l+1}(k)>P_{k^{*}}$ where $P_{k^{*}}^{l+1}(k)$ is the probability of selecting the best alternative through a pure greedy algorithm at round $l+1$ if $\mathbf{z}_{k}$ is tested at round $l$.
By definition

$$
P_{k^{*}}=\operatorname{Prob}\left\{\widehat{\boldsymbol{\mu}}_{k^{*}}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}>0\right\}
$$

and, under the assumption of Gaussianity,

$$
\left(\widehat{\boldsymbol{\mu}}_{k^{*}}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}\right) \sim \mathcal{N}\left(\mu_{k^{*}}-\mu_{\bar{k}^{*}}, \frac{\sigma_{k^{*}}^{2}}{n_{k^{*}}}+\frac{\sigma_{\bar{k}^{*}}^{2}}{n_{\bar{k}^{*}}}\right) .
$$

The mean of the Gaussian is positive. Reducing the variance of $\left(\widehat{\boldsymbol{\mu}}_{k^{*}}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}\right)$ will improve the $P_{k^{*}}$ at round $l+1$ (see fig. 1). Testing $\mathbf{z}_{k^{*}}$ or $\mathbf{z}_{\bar{k}^{*}}$ will both reduce the variance of $\left(\widehat{\boldsymbol{\mu}}_{k^{*}}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}\right)$, this proves the proposition.

Theorem 3. Let us consider a bandit problem with two normal arms $\mathbf{z}_{\#}=$ $\left\{\mathbf{z}_{k_{1}}, \mathbf{z}_{k_{2}}\right\}$. The following policy

$$
\left\{\begin{array}{l}
\text { if } N_{\Delta}<0 \text { then test } \mathbf{z}_{k_{1}} \\
\text { if } N_{\Delta}>0 \text { then test } \mathbf{z}_{k_{2}} \\
\text { if } N_{\Delta}=0 \text { then test either } \mathbf{z}_{k_{1}} \text { or } \mathbf{z}_{k_{2}}
\end{array}\right.
$$

will maximize $\mu_{g}$ at round $l+1$ where

$$
N_{\Delta}=n_{k_{1}} \cdot\left(n_{k_{1}}+1\right) \cdot\left(\sigma_{k_{2}}^{2}-\sigma_{k_{1}}^{2}\right)+\sigma_{k_{1}}^{2} \cdot\left(n_{k_{2}}+n_{k_{1}}+1\right) \cdot\left(n_{k_{1}}-n_{k_{2}}\right)
$$

Proof. If $\left[\mu_{g}^{l+1}\left(k_{1}\right)-\mu_{g}^{l+1}\left(k_{2}\right)\right]$ is positive then $\mathbf{z}_{1}$ must be tested to maximize $\mu_{g}$ at round $l+1$ and if $\left[\mu_{g}^{l+1}\left(k_{1}\right)-\mu_{g}^{l+1}\left(k_{2}\right)\right]$ is negative then $\mathbf{z}_{2}$ must be tested.


Fig. 1. If $K=2$, the probability of selecting the best random variable $P_{k^{*}}$ is the surface under the Gaussian when the abscissa takes his value in the set $[0,+\infty]$. Note that the mean of the Gaussian is always positive.

$$
\begin{aligned}
& \mu_{g}^{l+1}\left(k_{1}\right)-\mu_{g}^{l+1}\left(k_{2}\right) \\
& =P_{k^{*}}^{l+1}\left(k_{1}\right) \cdot \mu_{k^{*}}+P_{\bar{k}^{*}}^{l+1}\left(k_{1}\right) \cdot \mu_{\bar{k}^{*}}-P_{k^{*}}^{l+1}\left(k_{2}\right) \cdot \mu_{k^{*}}-P_{\bar{k}^{*}}^{l+1}\left(k_{2}\right) \cdot \mu_{\bar{k}^{*}} \\
& =\left[P_{k^{*}}^{l+1}\left(k_{1}\right)-P_{k^{*}}^{l+1}\left(k_{2}\right)\right] \cdot \mu_{k^{*}}+\left[P_{\bar{k}^{*}}^{l+1}\left(k_{1}\right)-P_{\bar{k}^{*}}^{l+1}\left(k_{2}\right)\right] \cdot \mu_{\bar{k}^{*}} \\
& =\left[P_{k^{*}}^{l+1}\left(k_{1}\right)-P_{k^{*}}^{l+1}\left(k_{2}\right)\right] \cdot \mu_{k^{*}}-\left[P_{k^{*}}^{l+1}\left(k_{1}\right)-P_{k^{*}}^{l+1}\left(k_{2}\right)\right] \cdot \mu_{\bar{k}^{*}} \\
& =\left[P_{k^{*}}^{l+1}\left(k_{1}\right)-P_{k^{*}}^{l+1}\left(k_{2}\right)\right] \cdot\left(\mu_{k^{*}}-\mu_{\bar{k}^{*}}\right) \\
& =\Delta P_{k^{*}} \cdot\left(\mu_{k^{*}}-\mu_{\bar{k}^{*}}\right)
\end{aligned}
$$

Since $\left(\mu_{k^{*}}-\mu_{\bar{k}^{*}}\right)>0$, the sign of $\left[\mu_{g}^{l+1}\left(k_{1}\right)-\mu_{g}^{l+1}\left(k_{2}\right)\right]$ is the same as the sign of $\Delta P_{k^{*}}$.

Let $\operatorname{Var}^{k_{1}}\left(\widehat{\boldsymbol{\mu}}_{k^{*}}^{l+1}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}^{l+1}\right)$ and $\operatorname{Var}^{k_{2}}\left(\widehat{\boldsymbol{\mu}}_{k^{*}}^{l+1}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}^{l+1}\right)$ be respectively the variance of $\left(\widehat{\boldsymbol{\mu}}_{k^{*}}^{l+1}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}^{l+1}\right)$ at $l+1$ if either $\mathbf{z}_{k_{1}}$ or $\mathbf{z}_{k_{2}}$ is tested at $l . \Delta \operatorname{Var}$ is defined as the difference between the two variances at $l+1$. A reduction of the variance of $\left(\widehat{\boldsymbol{\mu}}_{k^{*}}^{l+1}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}^{l+1}\right)$ improves the probability of selecting the best random variable
$\mathbf{z}_{k^{*}}$. The sign of $\Delta \operatorname{Var}$ is thus the opposite of the sign of $\Delta P_{k^{*}}$.

$$
\begin{aligned}
\Delta \operatorname{Var} & =\operatorname{Var}^{k_{1}}\left(\widehat{\boldsymbol{\mu}}_{k^{*}}^{l+1}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}^{l+1}\right)-\operatorname{Var}^{k_{2}}\left(\widehat{\boldsymbol{\mu}}_{k^{*}}^{l+1}-\widehat{\boldsymbol{\mu}}_{\bar{k}^{*}}^{l+1}\right) \\
& =\frac{\sigma_{k_{1}}^{2}}{n_{k_{1}}+1}+\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}-\frac{\sigma_{k_{1}}^{2}}{n_{k_{1}}}-\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}+1} \\
& =\frac{N_{\Delta}}{D_{\Delta}}
\end{aligned}
$$

where

$$
\begin{aligned}
N_{\Delta} & =n_{k_{2}} n_{k_{1}}\left(n_{k_{2}}+1\right) \sigma_{k_{1}}^{2}+\left(n_{k_{1}}+1\right) n_{k_{1}}\left(n_{k_{2}}+1\right) \sigma_{k_{2}}^{2} \\
& -\left(n_{k_{1}}+1\right) n_{k_{2}}\left(n_{k_{2}}+1\right) \sigma_{k_{1}}^{2}-\left(n_{k_{1}}+1\right) n_{k_{2}} n_{k_{1}} \sigma_{k_{2}}^{2} \\
& =n_{k_{1}}\left(\sigma_{k_{2}}^{2}-\sigma_{k_{1}}^{2}\right)\left(n_{k_{1}}+1\right)+\sigma_{k_{1}}^{2}\left(n_{k_{1}}-n_{k_{2}}\right)\left(n_{k_{1}}+n_{k_{2}}+1\right) \\
D_{\Delta} & =\left(n_{k_{1}}+1\right) n_{k_{2}} n_{k_{1}}\left(n_{k_{2}}+1\right)
\end{aligned}
$$

Since $D_{\Delta}>0$, the following policy

$$
\left\{\begin{array}{l}
\text { if } N_{\Delta}<0 \text { then test } \mathbf{z}_{k_{1}} \\
\text { if } N_{\Delta}>0 \text { then test } \mathbf{z}_{k_{2}} \\
\text { if } N_{\Delta}=0 \text { then test either } \mathbf{z}_{k_{1}} \text { or } \mathbf{z}_{k_{2}}
\end{array}\right.
$$

maximizes $\mu_{g}$ at round $l+1$.

Theorem 4. Let $\mathbf{z}_{\#}$ be a set containing two normal random variables $\mathbf{z}_{k_{1}}$ and $\mathbf{z}_{k_{2}}$. Suppose that the player adopts a bandit strategy which always tests the same $\operatorname{arm} \mathbf{z}_{k_{1}} \cdot \mathbf{z}_{k_{2}}$ is the other arm which is never tested. If

$$
\mathbf{X} \sim \mathcal{N}\left(\mu_{k_{1}}-\mu_{k_{2}}, \frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}\right)
$$

then

$$
\forall l \geq 1, \quad \mu_{g}<u b^{\infty}<\mu_{k^{*}}
$$

where

$$
u b^{\infty}=\left(\mu_{k_{1}}-\mu_{k_{2}}\right) \cdot \operatorname{Prob}\{\mathbf{X}>0\}+\mu_{k_{2}}
$$

Proof. Let $\tau$ be a positive integer, we can find an upper-bound for $\mu_{g}$ (see Theorem 2) :

$$
\begin{aligned}
\mu_{g} & <\mu_{g}^{l+1}\left(k_{1}\right) \\
& \leq \mu_{g}^{l+\tau}\left(k_{1}\right)
\end{aligned}
$$

The future expected greedy gain when arm $k_{1}$ is tested $\tau \geq 1$ times is $\mu_{g}^{l+\tau}\left(k_{1}\right)$ and is analytically defined as follows

$$
\begin{aligned}
\mu_{g}^{l+\tau}\left(k_{1}\right) & =\mu_{k_{1}} \cdot \operatorname{Prob}\left\{\left(\widehat{\boldsymbol{\mu}}_{k_{1}}-\widehat{\boldsymbol{\mu}}_{k_{2}}\right)_{\tau}>0\right\}+\mu_{k_{2}} \cdot\left(1-\operatorname{Prob}\left\{\left(\widehat{\boldsymbol{\mu}}_{k_{1}}-\widehat{\boldsymbol{\mu}}_{k_{2}}\right)_{\tau}>0\right\}\right) \\
& =\left(\mu_{k_{1}}-\mu_{k_{2}}\right) \cdot \operatorname{Prob}\left\{\left(\widehat{\boldsymbol{\mu}}_{k_{1}}-\widehat{\boldsymbol{\mu}}_{k_{2}}\right)_{\tau}>0\right\}+\mu_{k_{2}} \\
& =u b^{\tau}
\end{aligned}
$$

where

$$
\left(\widehat{\boldsymbol{\mu}}_{k_{1}}-\widehat{\boldsymbol{\mu}}_{k_{2}}\right)_{\tau} \sim \mathcal{N}\left(\mu_{k_{1}}-\mu_{k_{2}}, \operatorname{Var}_{\tau}\right)
$$

and where

$$
\operatorname{Var}_{\tau}=\frac{\sigma_{k_{1}}^{2}}{n_{k_{1}}+\tau}+\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}
$$

If the player tests infinitely the $\operatorname{arm} \mathbf{z}_{k_{1}}$ without reserving any round to explore $\mathbf{z}_{k_{2}}$ then the variance of $\left(\widehat{\boldsymbol{\mu}}_{k_{1}}-\widehat{\boldsymbol{\mu}}_{k_{2}}\right)_{\tau}$ converges

$$
\lim _{\tau \rightarrow \infty} \operatorname{Var}_{\tau}=\lim _{\tau \rightarrow \infty}\left(\frac{\sigma_{k_{1}}^{2}}{n_{k_{1}}+\tau}+\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}\right)=\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}
$$

and we define the random variable $\mathbf{X}$ as follows

$$
\mathbf{X}=\left(\widehat{\boldsymbol{\mu}}_{k_{1}}-\widehat{\boldsymbol{\mu}}_{k_{2}}\right)_{\infty} \sim \mathcal{N}\left(\mu_{k_{1}}-\mu_{k_{2}}, \frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}\right)
$$

The upper-bound of $\mu_{g}$ is defined as follows

$$
u b^{\infty}=\left(\mu_{k_{1}}-\mu_{k_{2}}\right) \cdot \operatorname{Prob}\{\mathbf{X}>0\}+\mu_{k_{2}}
$$

and we have $\mu_{g}<u b^{\infty}$. $u b^{\infty}$ reaches $\mu_{k^{*}}$ only if $P_{k^{*}}=1$ and $P_{k^{*}}$ equals one only if $\frac{\sigma_{k_{1}}^{2}}{n_{k_{1}}}+\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}=0$ (see figure 1). In the case of a bandit strategy which never tests $\mathbf{z}_{k_{2}}$, the variance $\frac{\sigma_{k_{1}}^{2}}{n_{k_{1}}}+\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}$ converges to $\frac{\sigma_{k_{2}}^{2}}{n_{k_{2}}}$ and we have thus $u b^{\infty}<\mu_{k^{*}}$.

Theorem 2 shows that, whatever the arm selected, the value of $\mu_{g}$ is always increasing. Theorem 3 proposes, knowing the characteristics of the arms, a new exploration strategy and Theorem 4 shows that the evolution of $\mu_{g}$ of any bandit strategy which never explores the arms and which always tests the same arm, is upper-bounded by $u b^{\infty}$ which is smaller than $\mu_{k^{*}}$. We will now empirically test these properties on six synthetic problems (table 1) with horizon $H=50$. Three algorithms are compared. At each round $l$, each of the three algorithms tests arm $\widetilde{k}$. The three algorithms are (i) an oracle greedy exploitation algorithm which always tests the best arm $\left[\widetilde{k}=k^{*}\right]$, (ii) a random exploration algorithm which periodically tests the two arms $[\widetilde{k}=(l \bmod 2)+1]$ and (iii) the optimal oracle exploration algorithm described in theorem 3 and called optiK2

|  | case 1 |  | case 2 |  | case 3 |  | case 4 |  | case 5 |  | case 6 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mu$ | $\sigma$ | $\mu$ | $\sigma$ | $\mu$ |  | $\mu$ | $\sigma$ | $\mu$ | $\sigma$ | $\mu$ |  |
| $\mathbf{z}_{1}$ | 0 | 2 |  | 1 | 0 | 3 | 0 | 2 | 0 | 0.01 | 0 | . 01 |
| $\mathbf{z}_{2}$ | 1 | 2 | 1 | 2 | 1 | 1 | 1 | 0.01 | 1 | 2 | 1 | 0.01 |

Table 1. Six synthetic problems to estimate the evolution of $\mu_{g}$ where $K=2$ under (i) an oracle greedy algorithm, (ii) a random exploration algorithm and (iii) the optiK2 algorithm. For each case, the first and the second column contain respectively the mean and the standard deviation of the random variables.
[if $N_{\Delta}<0$ then $\widetilde{k}=1$ else $\widetilde{k}=2$ ]. Initially, $n_{1}$ and $n_{2}$ are equal to one. At each round $l$ and following one of the three algorithms, the value of $n_{\tilde{k}}$ is increased by one and $\mu_{g}(l)$ (see eq. (2)) is computed. The evolution of $\mu_{g}$ in the time for the six synthetic problems are shown in figure 2 .

In all the cases, the curve of $\mu_{g}(l)$ is increasing (see Theorem 2) and optiK2 is optimal (see Theorem 3). In case $1, \sigma_{1}$ is equal to $\sigma_{2}$ such that $N_{\Delta}=\sigma_{1}^{2}\left(n_{1}-\right.$ $\left.n_{2}\right)\left(n_{1}+n_{2}+1\right)$ and optiK2 is then equivalent to a random exploration strategy. In case 5 , the standard deviation of the worst arm is very small such that the term $n_{1}\left(\sigma_{2}^{2}-\sigma_{1}^{2}\right)\left(n_{1}+1\right)$ in $N_{\Delta}$ is very high and optiK2 is equivalent to an oracle greedy strategy which always tests $\mathbf{z}_{2}$. Note that, with a random exploration strategy, we have

$$
\lim _{l \rightarrow \infty} \sqrt{\frac{\sigma_{k^{*}}^{2}}{n_{k^{*}}}+\frac{\sigma_{\bar{k}^{*}}^{2}}{n_{\bar{k}^{*}}}}=0
$$

and thus (see figure 1)

$$
\lim _{l \rightarrow \infty} P_{k^{*}}=1 .
$$

This means that playing a great number of random exploration actions guarantees that $\mu_{g}$ will always be improved and that, asymptotically, the value of $\mu_{g}$ will reach the value of $\mu_{k^{*}}$. In the case of bandit strategies which always test the same arm (like the oracle greedy strategy), we are guaranteed that $\mu_{g}$ will always be improved but $\mu_{g}$ is upper-bounded by $u b^{\infty}$ which is smaller than $\mu_{k^{*}}$ (see Theorem 4). Without exploration, the evolution of the gain of a greedy action is upper-bounded by $u b^{\infty}$ and the only way to increase this upper-bound is to explore the other arm.

### 3.2 The evolution of $\mu_{g}$ in the case where $K>2$

Consider now the case where $K>2$ and let $\mu_{g}^{l+1}(k)$ be the next greedy expected reward when arm $k$ is tested at round $l$. In this case, we will empirically show that the evolution of $\mu_{g}$ is more complex and that the monotone property $\mu_{g}^{l+1}(k)>$ $\mu_{g}$ shown in Theorem 2 does not always hold. We will consider the evolution of $\mu_{g}$ under a pure exploitation and a pure exploration policy.

First, consider the evolution of $\mu_{g}$ under a pure oracle exploitation greedy algorithm which knows the index $k^{*}$ and always tests this best arm. We will show that in some cases, mainly when the $\frac{\sigma}{\sqrt{n}}$ of the bad alternatives is high,


Fig. 2. Evolution of $\mu_{g}$ for the six synthetic problems under an oracle greedy algorithm, random exploration algorithm and the optiK2 algorithm. For each case, the value of the upper-bound $u b^{\infty}$ is given.
the value of $\mu_{g}^{l+1}(k)-\mu_{g}$ can be negative such that an oracle greedy action can reduce the value of $\mu_{g}$ at the next round $l+1$.

Six synthetic example cases (see table 2 and fig. 3) are tested in which $n_{k^{*}}$ takes its value in $[1,200]$ and where $\sigma_{k^{*}}$ is always equal to 10 . The first three cases concern problems where $K=3$ and in the last three cases $K=4$. In the cases 1,2 and 5 the $\frac{\sigma}{\sqrt{n}}$ of the worst arms has the same value, a small value in case 1 and in case 5 and a higher value in case 2 . The plots of the cases 1 and 5 show that increasing the $n_{k^{*}}$ of the best alternative, increases $\mu_{g}$ if the $\frac{\sigma}{\sqrt{n}}$ of the bad arms is small. Case 2 is the opposite : if the $\frac{\sigma}{\sqrt{n}}$ of the bad arms is high, then increasing the number of tests made on the best alternative, decreases $\mu_{g}$. In the cases 3,4 and 6 the $\frac{\sigma}{\sqrt{n}}$ of the worst alternatives is not the same and the corresponding plots in figure 3 are more non-linear. At the beginning (i.e. when the $\frac{\sigma_{k^{*}}}{\sqrt{n_{k^{*}}}}$ is high), testing the best $\mathbf{z}_{k^{*}}$ will degrade $\mu_{g}$ and when $\frac{\sigma_{k^{*}}}{\sqrt{n_{k^{*}}}}$ is small enough, testing $\mathbf{z}_{k^{*}}$ will improve $\mu_{g}$. Note that, like in case 4 , this improvement of $\mu_{g}$ can be very small. The two last cases are example cases where the means of the bad alternatives are different.

Note that in the cases $1,2,3$ and 4 , we have $\mu_{k^{*}}=1$ and the other means equal zero. In these four first cases, the expected gain of a greedy action $\mu_{g}$ equals thus the probability of correct selection $P_{k^{*}}$. In case 1, the $P_{k^{*}}$ evolution,

| Oracle exploitation greedy problem $K>2$ |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | case 1 |  | case 2 |  | case 3 |  | case 4 |  | case 5 |  | case 6 |  |
|  | $\mu$ | $\frac{\sigma}{\sqrt{n}}$ | $\mu$ | $\frac{\sigma}{\sqrt{n}}$ | $\mu$ | $\frac{\sigma}{\sqrt{n}}$ | $\mu$ | $\frac{\sigma}{\sqrt{n}}$ | $\mu$ | $\frac{\sigma}{\sqrt{n}}$ | $\mu$ | $\frac{\sigma}{\sqrt{n}}$ |
| $\mathbf{z}_{1}$ | 0 | 0.0001 | 0 | 5 | 0 | 1 | 0 | 0.0001 | 0 | 0.0001 | 0 | 0.0001 |
| $\mathbf{z}_{2}$ | 0 | 0.0001 | 0 | 5 | 0 | 10 | 0 | 4 | 0.2 | 0.0001 | 0.2 | 4 |
| $\mathbf{z}_{3}$ | 1 | . | 1 | - | 1 | . | 0 | 6 | 0.7 | 0.0001 | 0.7 | 6 |
| $\mathbf{z}_{4}$ |  |  |  |  |  |  | 1 | . | 1 | . | 1 | . |

Table 2. Six synthetic problems to estimate the evolution of $\mu_{g}$ under an oracle exploitation greedy algorithm where $K>2$. For each case, the first column contains the mean of the random variable and the second one contains the standard deviation of the sample average.
under an oracle greedy algorithm, is monotone increasing and in case $2, P_{k^{*}}$ is monotone decreasing. In the cases 3 and 4 , the $P_{k^{*}}$ evolution is more non-linear. Given a set of $K$ random variables $\left\{\mathbf{z}_{1}, \ldots, \mathbf{z}_{K}\right\}$, where $n_{k}$ is the number of observed rewards from $\mathbf{z}_{k}$ and where $\widehat{\boldsymbol{\mu}}_{k}$ is the corresponding sample average of $\mathbf{z}_{k}$. We have

$$
\begin{align*}
P_{k^{*}} & =\operatorname{Prob}\left\{\widehat{\boldsymbol{\mu}}_{k^{*}}>\widehat{\boldsymbol{\mu}}_{1}, \ldots, \widehat{\boldsymbol{\mu}}_{k^{*}}>\widehat{\boldsymbol{\mu}}_{k^{*}-1}, \widehat{\boldsymbol{\mu}}_{k^{*}}>\widehat{\boldsymbol{\mu}}_{k^{*}+1}, \ldots, \widehat{\boldsymbol{\mu}}_{k^{*}}>\widehat{\boldsymbol{\mu}}_{K}(3)\right. \\
& =\operatorname{Prob}\left\{\widehat{\boldsymbol{\mu}}_{k^{*}}>\mathbf{X}_{\max }\right\} \tag{4}
\end{align*}
$$

where

$$
\mathbf{X}_{\max }=\max _{k \in[1 K] / k^{*}} \widehat{\boldsymbol{\mu}}_{k}
$$

In equation 3 , the probability of correct selection is the probability that $\widehat{\boldsymbol{\mu}}_{k^{*}}$ is the highest in a set of $K$ random variables $\left\{\widehat{\boldsymbol{\mu}}_{1}, \ldots, \widehat{\boldsymbol{\mu}}_{K}\right\}$. In equation 4 , the bandit problem is transformed and the probability of correct selection is the probability that $\widehat{\boldsymbol{\mu}}_{k^{*}}$ is the highest in a set of only 2 random variables $\left\{\widehat{\boldsymbol{\mu}}_{k^{*}}, \mathbf{X}_{\max }\right\}$. Figure 4 gives the probability density function of $\mathbf{X}_{\max }$ in the four cases.

In case 1 , the $k^{*}$ th arm is confronted with $\mathbf{X}_{\max }$ which has a mean which is almost equal to zero. In case 1 , testing arm $k^{*}$ reduces the standard deviation of $\widehat{\boldsymbol{\mu}}_{k^{*}}$ and increases thus the probability that arm $k^{*}$ will be selected by a greedy action. In case 2 , the special arm $\mathbf{X}_{\max }$ has a higher mean than the mean of $\widehat{\boldsymbol{\mu}}_{k^{*}}$ which is one. $\widehat{\boldsymbol{\mu}}_{k^{*}}$ has more chance to win with a higher standard deviation. Testing arm $k^{*}$ reduces the standard deviation of $\widehat{\boldsymbol{\mu}}_{k^{*}}$ and decreases the probability that arm $k^{*}$ will be selected by a greedy action.

Note the non-regularity of the $\mathbf{X}_{\max }$ probability density functions of cases 3 and 4.

We now consider the evolution of $\mu_{g}$ when $K>2$ under a pure random exploration algorithm which periodically tests the two arms. Six synthetic example cases (see table 3 and fig. 5) are tested with horizon $H=200$. We observe that, in all six cases, the value of $\mu_{g}$ is always increasing when a random exploration algorithm is applied. The rate of the increase is high at the beginning and lower at the end.


Fig. 3. Evolution of $\mu_{g}$ of the six synthetic problems under an oracle exploitation greedy algorithm. Note that in the cases 1, 2, 3 and 4 , the expected gain of a greedy action $\mu_{g}$ equals the probability of correct selection $P_{k^{*}}$

As conclusion, applying a greedy policy when the $\frac{\sigma}{\sqrt{n}}$ of the bad arms is high decreases the value of the next $\mu_{g}$. It is better to make exploration at the beginning to improve $\mu_{g}$ and when the $\frac{\sigma}{\sqrt{n}}$ of the bad arms is small enough the player must move to a more greedy exploitation mode.

## 4 Conclusion

In this paper, we give an analytical definition of the expected gain of the wellknown greedy exploitation action and study its evolution in the context of the bandit problem. In the case where $K=2$, we give an optimal exploration strategy and we show that testing an arm will always improve the next greedy expected gain $\mu_{g}$. Results show also that, the $\mu_{g}$ evolution of strategies which always test the same arm are upper-bounded by a smaller value than $\mu_{k^{*}}$. In the case where $K>2$, the mode in which a semi-uniform bandit policy is working affects the evolution of $\mu_{g}$. At the beginning of the problem or when the reward variances of the bandit machines are high, a greedy action can reduce the next $\mu_{g}$ and random exploration is thus a better working mode. When enough knowledge is collected, such that the $\frac{\sigma}{\sqrt{n}}$ values are small, the policy must converge to a more greedy exploitation mode. Future work will focus on the estimation


Fig. 4. Monte Carlo determination of $\mathbf{X}_{\max }$ for the cases 1, 2, 3 and 4 . The mean and the standard deviation are estimated in all the cases.
of $\mu_{g}$ based on the historical samples $\left\{\mathbf{Z}_{k}(l)\right\}, k=1, \ldots, K$ and on the use of these $\mu_{g}$ estimators to propose a new bandit algorithm.

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| Random exploration problem $K>2$ |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | case 1 |  | case 2 |  | case 3 |  | case 4 |  | case 5 |  | case 6 |  |
|  | $\mu$ | $\sigma$ | $\mu$ | $\sigma$ | $\mu$ | $\sigma$ | $\mu$ | $\sigma$ | $\mu$ | $\sigma$ | $\mu$ | $\sigma$ |
| $\mathbf{z}_{1}$ | 0 | 1 | 0 | 3 | 0 | 3 | 0 | 3 | 0 | 6 | 0 | 6 |
| $\mathbf{z}_{2}$ | 0 | 1 | 0 | 1 | 0 | 3 | 0 | 0.1 | 0.2 | 3 | 0.2 | 3 |
| $\mathbf{Z}_{3}$ | 1 | 0.1 | 1 | 0.1 | 1 | 0.1 | 1 | 3 | 0.7 | 0.1 | 0.7 | 0.1 |
| $\mathbf{z}_{4}$ |  |  |  |  |  |  |  |  | 1 | 3 | 1 | 0.1 |

Table 3. Six synthetic problems to estimate the evolution of $\mu_{g}$ under a random exploration algorithm where $K>2$. For each case, the first and the second column contain respectively the mean and the standard deviation of the random variables.
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Fig. 5. Evolution of $\mu_{g}$ of the six synthetic problems under a random exploration algorithm.


[^0]:    ${ }^{1}$ We use boldface symbols to denote random variables.

